

On some invariants in numerical semigroups and estimations of the order bound.

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¹ **Abstract.** Let $S = \{s_i\}_{i \in \mathbb{N}} \subseteq \mathbb{N}$ be a numerical semigroup. For $s_i \in S$, let $\nu(s_i)$ denote the number of pairs $(s_i - s_j, s_j) \in S^2$. When S is the Weierstrass semigroup of a family $\{\mathcal{C}_i\}_{i \in \mathbb{N}}$ of one-point algebraic-geometric codes, a good bound for the minimum distance of the code \mathcal{C}_i is the Feng and Rao *order bound* $d_{ORD}(\mathcal{C}_i)$. It is well-known that there exists an integer m such that $d_{ORD}(\mathcal{C}_i) = \nu(s_{i+1})$ for each $i \geq m$. By way of some suitable parameters related to the semigroup S , we find upper bounds for m and we evaluate m exactly in many cases. Further we conjecture a lower bound for m and we prove it in several classes of semigroups.

Index Therms. Numerical semigroup, Weierstrass semigroup, AG code, order bound on the minimum distance, Cohen-Macaulay type.

1 Introduction.

Let $S = \{s_i\}_{i \in \mathbb{N}} \subseteq \mathbb{N}$ be a numerical semigroup and let e, c, c', d, d' denote respectively the multiplicity, the conductor, the subconductor, the dominant of the semigroup and the greatest element in S preceding c' (if $e > 1$), as in Setting 2.1. Further let ℓ be the number of gaps of S between d and c , and let

$$\tilde{s} := \max\{s \in S \text{ such that } s \leq d \text{ and } s - \ell \notin S\}.$$

When S is the Weierstrass semigroup of a family $\{\mathcal{C}_i\}_{i \in \mathbb{N}}$ of one-point AG codes (see [3],[2]), a good bound for the minimum distance of \mathcal{C}_i is the Feng and Rao *order bound*

$$d_{ORD}(\mathcal{C}_i) := \min\{\nu(s_j) : j \geq i + 1\}$$

where, for $s_j \in S$, $\nu(s_j)$ denotes the number of pairs $(s_j - s_k, s_k) \in S^2$. It is well-known that there exists an integer m such that sequence $\{\nu(s_i)\}_{i \in \mathbb{N}}$ is non-decreasing for $i \geq m + 1$ (see[7]) and so

$$d_{ORD}(\mathcal{C}_i) = \nu(s_{i+1}) \text{ for } i \geq m.$$

For this reason it is important to find the element s_m of S . In our papers [5] and [6], we proved that $s_m = \tilde{s} + d$ if $\tilde{s} \geq d'$, and we evaluated s_m in cases $\ell \leq 2$, or $e \leq 6$, or *Cohen-Macaulay type* ≤ 3 .

In this paper, by a more detailed study of the semigroup we find interesting relations among the integers defined above; further by using these relations we deduce the Feng and Rao order bound in several new situations. Moreover in every considered case we show that $s_m \geq c + d - e$.

In Section 2, we establish various formulas and inequalities among the integers e, ℓ, d', c', d, c and $t := d - \tilde{s}$, see in particular (2.5) and (2.6).

In Section 3, by using the results of Section 2 and some result from [6], we improve the known facts on s_m recalled above; further we state for each semigroup and we prove in many cases the

$$\text{conjecture: } s_m \geq c + d - e.$$

In Section 4 we treat some particular cases: for each of them we also prove that the conjecture holds.

In conclusion we see that the value of the *order bound* s_m depends essentially on the position of the

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integer \tilde{s} in the semigroup. We summarize below the main results for the convenience of the reader.

$$\begin{aligned} \text{If } \tilde{s} &\geq d' + c' - d && \text{then } s_m = \tilde{s} + d \\ \text{if } \tilde{s} &= 2d' - d && \text{then } s_m = \tilde{s} + d = 2d' \\ \text{if } \tilde{s} &\leq d' + c' - d - 1 && \text{then } s_m \leq \max\{\tilde{s} + d, 2d'\} \end{aligned} \quad (3.4).$$

When $\tilde{s} \leq d' + c' - d - 1$, we prove the following partial answers.

$$\text{If } [d' - \ell, d'] \cap \mathbb{N} \subseteq S \quad \text{then} \quad \begin{cases} \tilde{s} + d' - \ell + 1 \leq s_m \leq 2d' & \text{if } 2d' - d < \tilde{s} < d' + c' - d \\ s_m = \tilde{s} + d & \text{otherwise.} \end{cases} \quad (3.11)$$

$$\text{If } \begin{cases} \tilde{s} \leq d' - 2 \\ [\tilde{s} + 2, d'] \cap \mathbb{N} \subseteq S \\ 2d' - d < \tilde{s} < d' + c' - d \end{cases} \quad \text{then} \quad \begin{cases} s_m = \tilde{s} + d & \iff \tilde{s} + 1 \notin S \text{ and } c' = d \\ s_m \leq \tilde{s} + d - 1 & \text{otherwise.} \end{cases} \quad (3.8)$$

$$\text{If } \begin{cases} \tilde{s} \leq 2d' - d \\ [\tilde{s} + 2, d'] \cap \mathbb{N} \subseteq S \end{cases} \quad \text{then } s_m = \tilde{s} + d. \quad (3.8)$$

Finally we consider several particular subcases: if H denotes the subset of gaps of S inside the interval $[c - e, c' - 1]$ and τ is the Cohen-Macaulay type of S , we deduce the exact value or good estimations for s_m in the following situations.

$$\begin{aligned} \text{If } H &= \emptyset, && \text{then } s_m = \tilde{s} + d \quad (4.1) \\ \text{If } H &\text{ is a non empty interval,} && \text{then } s_m = \begin{cases} 2d' & \text{if } \tilde{s} \geq 2d' + 1 - d \\ \tilde{s} + d & \text{otherwise} \end{cases} \quad (4.1) \\ \text{If } S &\text{ is associated to a Suzuki curve,} && \text{then } s_m = \tilde{s} + d \quad (4.14). \end{aligned}$$

$$\text{If } \#H \leq 2, \quad \text{see } (4.4).$$

$$\text{If } \ell \leq 3, \quad \text{see } (4.5), (4.7).$$

$$\text{If } \tau \leq 7, \quad \text{see } (4.10).$$

$$\text{If } e \leq 8, \quad \text{see } (4.11).$$

$$\text{If } S \text{ is generated by an Almost Arithmetic Sequence and } \text{embdim}(S) \leq 5, \text{ see } (4.12).$$

2 Semigroups: invariants and relations.

We begin by giving the setting of the paper.

Setting 2.1 In all the article we shall use the following notation. Let \mathbb{N} denote the set of all non-negative integers. A *numerical semigroup* is a subset S of \mathbb{N} containing 0, closed under summation and with finite complement in \mathbb{N} ; we shall always assume $S \neq \mathbb{N}$. We denote the elements of S by $\{s_i\}$, $i \in \mathbb{N}$, with $s_0 = 0 < s_1 < \dots < s_i < s_{i+1} \dots$.

We list below some invariants and subsets related to a semigroup $S \subset \mathbb{N}$ we shall need in the sequel.

$$\begin{aligned}
e &:= s_1 > 1, \text{ the multiplicity of } S. \\
c &:= \min \{r \in S \mid r + \mathbb{N} \subseteq S\}, \text{ the conductor of } S \\
d &:= \text{the greatest element in } S \text{ preceding } c, \text{ the dominant of } S \\
c' &:= \max\{s_i \in S \mid s_i \leq d \text{ and } s_i - 1 \notin S\}, \text{ the subconductor of } S \\
d' &:= \text{the greatest element in } S \text{ preceding } c', \text{ when } c' > 0 \\
\ell &:= c - 1 - d, \text{ the number of gaps of } S \text{ greater than } d \\
g &:= \#(\mathbb{N} \setminus S), \text{ the genus of } S \text{ (= the number of gaps of } S) \\
S' &:= \{s \in S \mid s \leq d'\} \subseteq S \\
S(1) &:= \{b \in \mathbb{N} \mid b + (S \setminus \{0\}) \subseteq S\} \\
\tau &:= \#(S(1) \setminus S), \text{ the Cohen–Macaulay type of } S \\
H &:= [c - e, c'] \cap \mathbb{N} \setminus S \subseteq \mathbb{N} \setminus S.
\end{aligned}$$

(Note that $c - e \leq c'$ since $c - e - 1 \notin S$).

According to this notation we can represent a semigroup S with $c' > 0$ as follows:

$$S = \{0, * \dots *, e, \dots, d', * \dots *, c' \xleftrightarrow{\ell \text{ gaps}} d, * \dots *, c \rightarrow\} = S' \cup \{c' \xleftrightarrow{\ell \text{ gaps}} d, * \dots *, c \rightarrow\},$$

where “ $*$ ” indicates *gaps*, “ $* \dots *$ ” *interval of all gaps*, and “ $\xleftrightarrow{\ell \text{ gaps}}$ ” *interval without gaps*.

Moreover for $s_i \in S$ we fix the following notation.

$$\begin{aligned}
N(s_i) &:= \{(s_j, s_k) \in S^2 \mid s_i = s_j + s_k\}; & \nu(s_i) &:= \#N(s_i); \\
\eta(s_i) &:= \nu(s_{i+1}) - \nu(s_i). \\
d_{ORD}(i) &:= \min\{\nu(s_j) \mid j > i\}, \text{ the order bound.} \\
A(s_i) &:= \{(x, y), (y, x) \in N(s_i) \mid x \leq d', c' \leq y \leq d\}; & \alpha(s_i) &:= \#A(s_{i+1}) - \#A(s_i). \\
B(s_i) &:= \{(x, y) \in [c', d]^2 \cap N(s_i)\}; & \beta(s_i) &:= \#B(s_{i+1}) - \#B(s_i). \\
C(s_i) &:= \{(x, y) \in S'^2 \cap N(s_i)\}; & \gamma(s_i) &:= \#C(s_{i+1}) - \#C(s_i). \\
D(s_i) &:= \{(x, y), (y, x) \in N(s_i) \mid x \geq c, x \geq y\}; & \delta(s_i) &:= \#D(s_{i+1}) - \#D(s_i).
\end{aligned}$$

Now we recall some definition and former results for completeness. First, a semigroup S is called

ordinary if $S = \{0\} \cup \{n \in \mathbb{N}, n \geq c\}$

acute if either S is ordinary, or c, d, c', d' satisfy $c - d \leq c' - d'$ [1, Def. 5.6].

Definition 2.2 We define the invariants \tilde{s} , m and t as follows.

$$\tilde{s} := \max \{s \in S \text{ such that } s \leq d \text{ and } s - \ell \notin S\}.$$

$$t := d - \tilde{s}.$$

$$m := \min \{j \in \mathbb{N} \text{ such that the sequence } \{\nu(s_i)\}_{i \in \mathbb{N}} \text{ is non-decreasing for } i > j\}$$

$$(m > 0 \iff \nu(s_m) > \nu(s_{m+1}) \text{ and } \nu(s_{m+k}) \leq \nu(s_{m+k+1}), \text{ for each } k \geq 1).$$

Theorem 2.3 Let $S = \{s_i\}_{i \in \mathbb{N}}$ be as in Setting 2.1.

$$(1) \nu(s_i) = i + 1 - g, \text{ for every } s_i \geq 2c - 1. \quad [7, \text{Th. 3.8}]$$

$$(2) \nu(s_{i+1}) \geq \nu(s_i), \text{ for every } s_i \geq 2d + 1. \quad [5, \text{Prop. 3.9.1}]$$

$$(3) \text{ If } S \text{ is an ordinary semigroup, then } m = 0. \quad [1, \text{Th. 7.3}]$$

$$(4) \text{ If } \tilde{s} \geq d', \text{ then } s_m = \tilde{s} + d \quad [6, \text{Th. 4.1, Th.4.2}].$$

In particular:

$$(a) \text{ if } t \leq 2, \text{ then } s_m = \tilde{s} + d,$$

$$(b) \text{ if } S \text{ is an acute semigroup, then } s_m = \tilde{s} + d, \text{ with}$$

- (i) either $d - c' \geq \ell - 1$, $s_m = c + c' - 2 = \tilde{s} + d$,
- (ii) or $\tilde{s} = d$ ($s_m = 2d$). [5, Prop. 3.4] .

(5) If $c' \in \{c - e, c - e + 1\}$, then S is acute. [6, Lemma 5.1].

Remark 2.4 (1) By the definition of \tilde{s} it is clear that:

$$s - \ell \in S \text{ for each } s \in S \text{ such that } \tilde{s} < s \leq d.$$

(2) Theorem 2.3 implies that $0 < s_m \leq 2d$ for every non-ordinary semigroup.

(3) The condition (a) of (2.3.4) does not imply S acute; analogously there exist non-acute semigroups satisfying the conditions (4.b, i - ii), see Example 2.9.2.

We complete this section with some general relation among the invariants defined above.

Proposition 2.5 [6, Prop. 2.5] Let $c' = c - e + q$, with $q \geq 0$. Then

(1) $e \leq 2\ell + t + q$.

(2) The following conditions

(a) $d - c' \geq \ell - 1$ (i.e. $c + c' - 2 \leq 2d$).

(b) $\tilde{s} - \ell = c' - 1$.

(c) $c + c' - 2 = \tilde{s} + d$.

(d) $e = 2\ell + t + q$

are equivalent and imply

(i) $c' \leq \tilde{s} \leq d$ ($\implies s_m = \tilde{s} + d$).

(ii) S is acute $\iff d - d' \geq 2\ell + t$.

Proof. (1) By definition 2.2 we have $\tilde{s} - \ell \leq c' - 1 = c - e + q - 1$, then $\tilde{s} - \ell \leq d + \ell - e + q$ and so $e \leq 2\ell + t + q$.

(2) The equivalences (2.a) \iff (2.b) \iff (2.c) are proved in [6, Prop. 2.5]. Clearly the equality $e = 2\ell + t + q$ holds if and only if $\tilde{s} - \ell = d - t - \ell = c' - 1$. Further:

(i) is obvious by (2.b).

(ii) If (2.b) holds, then $d - d' - (2\ell + t) = \tilde{s} - \ell - d' - \ell = (c' - d') - (\ell + 1) = (c' - d') - (c - d)$. Then S is acute $\iff d - d' \geq 2\ell + t$.

Theorem 2.6 Let $t = d - \tilde{s}$ (see 2.1). The following facts hold.

(1) (a) If $0 \leq h < e$ and $d - h \in S$, then $e \geq h + \ell + 1$.

(b) If $s, s' \in S$, $s \geq c - e$ and $s - \ell \leq s' < s$, then $s' \geq c - e$.

(2) $\tilde{s} \geq c - e$ (equivalently, $e \geq t + \ell + 1$).

(3) Let $t > 0$ and let $s' := \min\{s \in S \mid s > \tilde{s}\}$. Then

$$e \geq 2\ell + 1 + d - s' \geq 2\ell + 1 \text{ (equivalently, } s' \geq c - e + \ell).$$

In particular,

(a) $\tilde{s} + 1 \in S \implies e \geq 2\ell + t$;

(b) $c' \leq \tilde{s} < d \implies e \geq 2\ell + t$.

(4) One of the following conditions hold

(a) $\tilde{s} - \ell \geq c - e$ (equivalently $e > 2\ell + t$, equivalently $\tilde{s} - \ell \in H$)

(b) $\tilde{s} - \ell = c - e - 1$ (equivalently $e = 2\ell + t$)

(c) $c - e - \ell \leq \tilde{s} - \ell < c - e - 1$ (equivalently $e < 2\ell + t$)

(5) Assume $e < 2\ell + t$, then :

- (a) either $\tilde{s} \leq d'$ or $t = 0$;
- (b) in case $\tilde{s} \leq d'$ we have: $[\tilde{s} + 1, c - e + \ell - 1] \cap S = \emptyset$, $\#H \geq 2\ell + t - e$, further if $\tilde{s} < d'$, then $\#H \geq 2\ell + t - e + 1 \geq 2$.

Proof. (1.a) We have $d < d - h + e \in S$. Hence $d - h + e \geq c = d + \ell + 1$.

(1.b) If $s \geq c$ we have $s' \geq c$. If $s \leq d$, let $s' = d - h$, $s = d - k \geq c - e$ (hence $k + \ell \leq e - 1$), then $d - \ell - k \leq d - h \implies h \leq k + \ell \leq e - 1$. Now apply (a): $e \geq h + \ell + 1$, equivalently, $d - h \geq d + \ell + 1 - e$.

(2) Let $d = c - e + h\ell + r$, with $h \geq 0$, $0 \leq r < \ell$ (recall that $c - e \leq d$). If $\tilde{s} < c - e$, by (2.4.1) we get $\tilde{s} < d - h\ell \in S$, $d - (h + 1)\ell \in S$; further we get $c - e - \ell \leq d - (h + 1)\ell < c - e$, a contradiction because $[c - \ell - e, c - e - 1] \cap S = \emptyset$ for every semigroup.

(3) By (2.4.1), $\tilde{s} < s' \leq d \implies s' - \ell \in S$ and so $s' - \ell + e \in S$. Since $c - e \leq \tilde{s} < s'$, we get $s' - \ell + e > c - \ell = d + 1$; it follows that $s' - \ell + e \geq c = d + \ell + 1$.

(4) Since $c - e - 1 \notin S$, the statements are almost immediate by (2).

(5.a) follows by ((3.b)).

(5.b) We are in case (4.c): since obviously $[c - e - \ell, c - e - 1] \cap S = \emptyset$, we have

$$[\tilde{s} - \ell + 1, c - e - 1] \cap \mathbb{N} \subseteq [c - \ell - e + 1, c - e - 1] \cap \mathbb{N} \subseteq \mathbb{N} \setminus S.$$

By the definition of \tilde{s} , we deduce that $[\tilde{s} + 1, c - e + \ell - 1] \cap \mathbb{N} \subseteq H$. The last inequality follows recalling that $\tilde{s} + 1 < d' \implies c - e + \ell - 1 < d'$, hence $d' + 1 \in H \setminus [\tilde{s} + 1, c - e + \ell - 1]$. \diamond

Corollary 2.7 Assume $\tilde{s} < d$ (i.e. $t > 0$). Then

- (1) If $c' = c - e + q$, with $q \in \{0, 1\}$, then $d - c' \geq \ell - 1$ and $e = 2\ell + t + q$.
- (2) If $d - c' \leq \ell - 2$, then
 - (a) $d - c' + 2 \leq d - d' \leq \ell \leq e - 3 - (d - c')$;
 - (b) if $\tilde{s} \geq 2d' - d$, then $t \leq 2\ell$.

Proof. (1) By the assumptions and by (5), (4) of Theorem 2.3, we have $d - c' \geq \ell - 1$. Then the other statement follows by (2.5.2).

(2.a) Since $d' \leq c' - 2$ the first inequality holds for any semigroup. We have $d - c' \leq \ell - 2$, by assumption, and $d - \ell \in S$, by (2.4.1). Hence $d' \geq d - \ell$. For the last inequality see [6, Prop. 5.2].

(2.b) follows by (2.a) because the assumption means $t \leq 2d - 2d'$. \diamond

Corollary 2.8 Assume $\tilde{s} \leq d'$. The following facts hold:

- (1) $d - c' \leq \ell - 2$, $d - d' \leq \ell$, $c' \geq c - e + 2$.
- (2) If $e \leq 2\ell + t$, then
 - (a) $\#H \leq \ell + t - 2(d - c') - 4$.
 - (b) $e = 2\ell + t \iff d + 2\ell - e \in S$.
- (3) If $H \subseteq [d' - t + 1, c' - 1]$, then $e \leq 2\ell + t$.
- (4) If $H = [d' + 1, c' - 1] \cap \mathbb{N}$ and $e < 2\ell + t$, then $\tilde{s} = d'$.

Proof. (1) By (2.5.2) we see that $\tilde{s} \leq d' \implies d - c' \leq \ell - 2$, therefore $d - d' \leq \ell$ (2.7). Further we have $c' \geq c - e + 2$ because $c - e \leq \tilde{s}$ (2.6.2) and $\tilde{s} \leq d' \leq c' - 2$.

(2.a) By (1) we have $d - d' \leq \ell$, then by (2.6.1-2) and by (2.4.1), we deduce that

$$\{c' - \ell, \dots, d - \ell, \tilde{s}, d'\} \subset S \cap [c - e, d'].$$

Hence $\#H \leq c' - (c - e) - 2 - (d - c' + 1) = 2c' - 2d - \ell - 1 - 3 + e \leq 2c' - 2d - \ell - 1 - 3 + 2\ell + t = \ell + t - 2(d - c') - 4$.

(2b) Clearly $e = 2\ell + t \implies d + 2\ell - e = \tilde{s} \in S$. The converse follows by the assumption and by Theorem 2.6.5b: $c - e + \ell - 1 = d + 2\ell - e \in S \implies e \geq 2\ell + t$; then $e = 2\ell + t$.

(3) $\tilde{s} \leq d' \implies d - \ell \leq d'$ by (1). Hence $\tilde{s} - \ell = d - \ell - t \leq d' - t$: now the assumption on H implies $\tilde{s} - \ell \notin H$, i.e., $e \leq 2\ell + t$ (2.6.4).

(4) If $e < 2\ell + t$, we have $\tilde{s} + 1 \notin S$ (2.6.3); since $c - e < \tilde{s} + 1$ we get $\tilde{s} + 1 \in H$, and so $\tilde{s} = d'$. \diamond

Example 2.9 (1) If $t > 0$, for each s_i such that $\tilde{s} < s_i \leq d$, we have that $s_i - \ell \in S$, hence $s_i - s_{i-1} \leq \ell$, but it is not true that for each $s_i \in S$ such that $c - e \leq s_i < d$, we have $s_{i+1} - s_i \leq \ell$: for instance let $\ell \geq 2$ and $S = \{0, 5\ell_e, 7\ell_{\tilde{s}=d-\ell=d'}, 8\ell_d, 9\ell + 1_c \rightarrow\}$.

(2) When $t = 0$ the inequality $e \geq 2\ell + 1$ (proved in (2.6.3) for $t > 0$) in general is not true, even for acute semigroups:

$$S_1 = \{0, 10_{e=d'}, 17_{c'}, 18, 19, 20_d, 27_c \rightarrow\} :$$

$$\ell = 6, t = 0, \tilde{s} = d \text{ and } S_1 \text{ is acute with } d - c' \leq \ell - 2, e < 2\ell.$$

$$S_2 = \{0, 8_e, 12_{d'}, 14_{c'}, 15, 16_d, 20_c \rightarrow\} :$$

$$\ell = 3, t = 0, S_2 \text{ is non-acute with } d - c' = \ell - 1, e > 2\ell.$$

$$S_3 = \{0, 7_e, 12_{d'}, 14_{c'=d}, 19_c \rightarrow\} :$$

$$\ell = 4, t = 0, S_3 \text{ is non-acute with } d - c' \leq \ell - 2, e < 2\ell.$$

$$S_4 = \{0, 10_{e=d'}, 14_d, 20_c \rightarrow\}$$

$$\ell = 5, t = 0, S_4 \text{ is non-acute with } d - c' \leq \ell - 2, e = 2\ell.$$

(3) When $\tilde{s} \leq d'$ we can have every case (a), (b), (c) of (2.6.4):

$$S_5 = \{0, 13_e, 15_{d'}, 20_d, 26_c \rightarrow\} : \ell = t = 5 \quad e < 2\ell + t = 15;$$

$$S_6 = \{0, 15_e, 19_{d'=\tilde{s}}, 24_d, 30_c \rightarrow\} : \ell = t = 5 \quad e = 2\ell + t = 15;$$

$$S_7 = \{0, 26_e, 28, 31_{d'}, 33_d, 39_c \rightarrow\} : \ell = t = 5 \quad e > 2\ell + t = 15.$$

3 General results on s_m .

As seen in [6], $s_m = \tilde{s} + d$ when $\tilde{s} \geq d'$. To give estimations of s_m in the remaining cases we shall use the same tools as in [6]: we recall them for the convenience of the reader. We shall add some improvement, as the general inequalities (3.1.3) on the difference $\nu(s+1) - \nu(s)$, however a great part of the following (3.1),(3.3),(3.4) is already proved in [6, 4.1, 4.2, 4.3].

Proposition 3.1 *Let $S' = \{s \in S, | s \leq d'\}$. For $s_i \in S$, let $\eta(s_i)$, $\alpha(s_i)$, $\beta(s_i)$, $\gamma(s_i)$, $\delta(s_i)$ be as in (2.1). Then:*

(1) *If $\tilde{s} < d'$, we have:*

$$s_{i+1} = s_i + 1 \text{ for } s_i \geq \tilde{s} + d' - \ell, \text{ in particular for } s_i \geq 2d'.$$

(2) *For each $s_i \in S$: $\eta(s_i) = \alpha(s_i) + \beta(s_i) + \gamma(s_i) + \delta(s_i)$. Further*

$$\alpha(s_i) = \begin{cases} -2 & \text{if } (s_{i+1} - c' \notin S' \text{ and } s_i - d \in S') \\ 0 & \text{if } (s_{i+1} - c' \in S' \iff s_i - d \in S') \\ 2 & \text{if } (s_{i+1} - c' \in S' \text{ and } s_i - d \notin S'). \end{cases}$$

$$\beta(s_i) = \begin{cases} 0 & \text{if } s_i \leq 2c' - 2 \text{ or } s_i > 2d \\ 1 & \text{if } 2c' - 1 \leq s_i \leq c' + d - 1 \\ -1 & \text{if } c' + d \leq s_i \leq 2d. \end{cases}$$

$$\gamma(s_i) = \begin{cases} 0 & \text{if } s_i \geq 2d' + 1 \\ -1 & \text{if } s_i = 2d' \\ -1 & \text{if } s_i < 2d' \text{ and } [s_i - d', d'] \cap \mathbb{N} \subseteq S. \end{cases}$$

$$\delta(s_i) = \begin{cases} 0 & \text{if } s_{i+1} - c \notin S, \quad s_i \leq 2c - 1 \\ 2 & \text{if } s_{i+1} - c \in S, \quad s_i \leq 2c - 1 \\ 1 & \text{if } s_i \geq 2c. \end{cases}$$

(3) *Let $s = 2d - k < 2d$ and $s + 1 \in S$, then:*

$$(a) -\left\lceil \frac{k}{2} \right\rceil - 1 \leq \nu(s+1) - \nu(s) \leq \left\lceil \frac{k+5}{2} \right\rceil.$$

$$(b) \text{ If } s = 2d' - h < 2d', \text{ then } -\left\lceil \frac{h}{2} \right\rceil - 1 \leq \gamma(s) \leq \left\lceil \frac{h+1}{2} \right\rceil.$$

Proof. (1) By assumption and by (2.6.2) we have $c - e \leq \tilde{s} < d'$ and so $d' - \ell \in S$. It follows $d' - \ell \geq e$ because $d' \geq e \geq \ell + t + 1$ (2.6.2). Hence $s \geq \tilde{s} + d' - \ell \implies s \geq c$.

(2) By [6, (3.3)...(3.7)] we have only to prove the last two statements for γ .

Let $s = 2d' - h \in S$, $h \in \mathbb{N}$, $s + 1 = 2d' - h + 1$ and assume $[d' - h, d'] = [s_i - d', d'] \cap \mathbb{N} \subseteq S$. Then:

$$\begin{aligned} C(s_i) &= \{(d' - h, d'), (d' - h + 1, d' - 1), (d' - h + 2, d' - 2), \dots, (d' - 1, d' - h + 1), (d', d' - h)\} \\ C(s_{i+1}) &= \{(d' - h + 1, d'), (d' - h + 2, d' - 1), \dots, (d' - 1, d' - h + 2), (d', d' - h + 1)\} \end{aligned}$$

it follows that $\gamma(s_i) = \#C(s_{i+1}) - \#C(s_i) = h - (h + 1) = -1$.

(3.a) To prove the inequalities for $s = 2d - k$, let $2d - k = x + y$, with $x \geq y$: then $\begin{cases} y \leq d - \left\lceil \frac{k}{2} \right\rceil \\ x \geq d - \left\lceil \frac{k}{2} \right\rceil \end{cases}$.

Therefore divide the interval $[d - \left\lceil \frac{k}{2} \right\rceil, d] \cap \mathbb{N}$ in subsets

$$\Lambda_j = [*** a_j \longleftrightarrow b_j] = H_j \cup S_j, \quad j = 1, \dots, j(s)$$

with $S_j \subseteq S$, $S_j = [a_j, b_j] \cap \mathbb{N}$ interval such that $b_j + 1 \notin S$, and $H_j \subseteq \mathbb{N} \setminus S$, $H_j \neq \emptyset$, if $j > 1$ (i.e. $a_{j-1} \notin S$ for $j > 1$, $H_1 = \emptyset \iff a_1 = d - \lfloor k/2 \rfloor \in S$).

Let $N_j(s) := N(s) \cap \{(x, y), (y, x) \mid x \in S_j, x \geq y\}$: we have $N(s) = \bigcup_j N_j(s) \cup D(s)$. Hence:

$$\nu(s+1) - \nu(s) = (\sum_j n_j) + \delta(s), \quad \text{where } n_j = \#N_j(s+1) - \#N_j(s).$$

Further: $-2 \leq n_j \leq 2$. This fact follows by the same argument used to prove the formulas for $\alpha(s_i)$, $\beta(s_i)$ recalled in statement (2) above. Since $0 \leq \delta(s) \leq 2$ (see (2) above) we conclude that

$$(*) \quad -2j(s) \leq \nu(s+1) - \nu(s) \leq 2j(s) + 2.$$

More precisely, to evaluate the largest and lowest possible values of $\nu(s+1) - \nu(s)$, with $s = 2d - k$,

we consider separately four cases: $\begin{cases} (A) & k = 4p \\ (B) & k = 4p + 1 \\ (C) & k = 4p + 2 \\ (D) & k = 4p + 3. \end{cases}$

In each case we can see that $j(s) \leq p + 1 = \left\lceil \frac{k}{4} \right\rceil + 1$. First note that $d \in S$, hence $j(s)$ is maximal when $\#H_j = \#S_j = 1$, i.e. $[d - \left\lceil \frac{k}{2} \right\rceil, d] = [\dots * \times * \times \dots * d]$ (where \times means element in S).

In each of the above cases we shall find integers x_1, x_2, y_1, y_2 such that $\begin{cases} x_1 \leq \nu(s+1) \leq x_2 \\ y_1 \leq \nu(s) \leq y_2 \end{cases}$, then the statement will follow by the obvious inequality $x_1 - y_2 \leq \nu(s+1) - \nu(s) \leq x_2 - y_1$.

- If either $k = 4p$, or $k = 4p + 1$, then $j(s)$ is maximal if and only if

$$[d - \left\lceil \frac{k}{2} \right\rceil, d] = [d - 2p * \dots * d - 2 * d], \quad \text{with } j(s) = p + 1.$$

Note that when $j(s) = p + 1$, then $1 \leq \#(N(s) \setminus D(s)) \leq 2p + 1$ because $(d - 2p, d - 2p) \in N(s)$; further we have $p \leq j(s+1) \leq p + 1$ and so $0 \leq \#N(s+1) \leq 2p + 4$.

If $k = 4p$, we have $1 \leq \#(N(s) \setminus D(s)) \leq 2p + 1$, since $(d - 2p, d - 2p) \in N(s)$, further $j(s+1) = p$, hence $0 \leq \#(N(s+1) \setminus D(s+1)) \leq 2p$.

$$-\left\lceil \frac{k}{2} \right\rceil - 1 = -2p - 1 \leq \nu(s+1) - \nu(s) \leq 2p + 2 - 1 < \left\lceil \frac{k+5}{2} \right\rceil.$$

If $k = 4p + 1$, we have $0 \leq \#(N(s) \setminus D(s)) \leq 2p + 2$, further $s + 1 = 2d - 4p$, therefore $1 \leq \#(N(s+1) \setminus D(s+1)) \leq 2p + 1$. We obtain:

$$-\left\lceil \frac{k}{2} \right\rceil - 1 = -2p - 1 \leq \nu(s+1) - \nu(s) \leq 2p + 3 = \left\lceil \frac{k+5}{2} \right\rceil.$$

- If either $k = 4p + 2$, or $k = 4p + 3$, then $\left\lfloor \frac{k}{2} \right\rfloor = 2p + 1$,
analogously we get $j(s) = p + 1$ maximal if

$$\left[d - \left\lfloor \frac{k}{2} \right\rfloor, d \right] = \left[\begin{array}{l} [*d - 2p * \dots * d - 2 * d] \text{ or} \\ [d - 2p - 1 \dots * \times \times * \dots d] \end{array} \right. \left. \left(\text{with one and only one } j_0 \text{ such that } \#S_{j_0} = 2 \right) \right].$$

If $k = 4p + 2$, in the first subcase we get $0 \leq (N(s) \setminus D(s)) \leq 2p + 2$, and $0 \leq \#(N(s+1) \setminus D(s+1)) \leq 2p$ because $(d - 2p - 1, d - 2p) \notin N(s + 1)$. Hence

$$-\left\lfloor \frac{k}{2} \right\rfloor - 1 = -2p - 2 \leq \nu(s + 1) - \nu(s) \leq 2p + 2 < \left\lfloor \frac{k + 5}{2} \right\rfloor.$$

In the second subcase we get $1 \leq \#(N(s) \setminus D(s)) \leq 2p + 2$, because $(d - 2p - 1, d - 2p - 1) \in N(s)$ and $0 \leq (N(s + 1) \setminus D(s + 1)) \leq 2p$ since $(d - 2p - 1, d - 2p) \notin N(s + 1)$. Hence

$$-\left\lfloor \frac{k}{2} \right\rfloor - 1 = -2p - 2 \leq \nu(s + 1) - \nu(s) \leq 2p + 1 < \left\lfloor \frac{k + 5}{2} \right\rfloor.$$

If $k = 4p + 3$, in the first subcase we get $0 \leq (N(s) \setminus D(s)) \leq 2p + 2$, and $0 \leq \#(N(s + 1) \setminus D(s + 1)) \leq 2p + 2$ because $(d - 2p - 1, d - 2p) \notin N(s + 1)$. Hence

$$-\left\lfloor \frac{k}{2} \right\rfloor - 1 = -2p - 2 \leq \nu(s + 1) - \nu(s) \leq 2p + 4 = \left\lfloor \frac{k + 5}{2} \right\rfloor.$$

In the second subcase we get $0 \leq \#(N(s) \setminus D(s)) \leq 2p + 2$ and $0 \leq \#(N(s + 1) \setminus D(s)) \leq 2p + 1$ because $(d - 2p - 1, d - 2p - 1) \in N(s + 1)$. Hence

$$-\left\lfloor \frac{k}{2} \right\rfloor - 1 = -2p - 2 \leq \nu(s + 1) - \nu(s) \leq 2p + 3 < \left\lfloor \frac{k + 5}{2} \right\rfloor.$$

(3.b) The proof is quite similar to the above one: since $\gamma(s) = \#C(s + 1) - \#C(s)$, we do not need to add $\delta(s)$ and so formula (*) becomes

$$-2j'(s) \leq \gamma(s) \leq 2j'(s),$$

where $j'(s)$ is the number of the subsets Λ_j as in (3.a) contained in the interval $[d' - \left\lfloor \frac{h}{2} \right\rfloor, d'] \cap \mathbb{N}$.

Then it suffices to proceed as above. \diamond

Example 3.2 The bounds found in (3.1.3a) are both sharp. To see this fact, consider $S = \{0, 10_e, 20_{d'}, 30_d, 40_c \rightarrow\}$ and the elements $s = 2d - 1 = 59$, $s + 1 = 2d = 60$. By a direct computation we easily get: $\nu(s + 1) - \nu(s) = 3 = \left\lfloor \frac{k + 5}{2} \right\rfloor$ (with $k = 1$), and

$$\nu(s + 2) - \nu(s + 1) = -\left\lfloor \frac{k}{2} \right\rfloor - 1 \quad (\text{with } k = 0).$$

Proposition 3.3 Let \circ mean $\notin S'$ and \times mean $\in S'$ (recall that for $s \leq d'$, we have $s \in S \iff s \in S'$). The following tables describe the difference $\eta(s_i) = \nu(s_{i+1}) - \nu(s_i)$ for $s_i \in S$, $s_i < 2c$ in function of $\alpha, \beta, \gamma, \delta$.

(a) If $s_i < 2c$:

$$\left[\begin{array}{cccccc} s_{i+1} - c & s_i - d & s_{i+1} - c' & \alpha & \delta & \eta(s_i) \\ \notin S & \times & \circ & -2 & 0 & \beta + \gamma - 2 \\ \notin S & \circ & \circ & 0 & 0 & \beta + \gamma \\ \notin S & \times & \times & 0 & 0 & \beta + \gamma \\ \notin S & \circ & \times & 2 & 0 & \beta + \gamma + 2 \\ \in S & \times & \circ & -2 & 2 & \beta + \gamma \\ \in S & \circ & \circ & 0 & 2 & \beta + \gamma + 2 \\ \in S & \times & \times & 0 & 2 & \beta + \gamma + 2 \\ \in S & \circ & \times & 2 & 2 & \beta + \gamma + 4 \end{array} \right].$$

More precisely we have the following subcases.

(b) If $s_i \leq 2d' - 1$, then $\beta = 0$:

$$\begin{bmatrix} s_{i+1} - c & s_i - d & s_{i+1} - c' & \alpha & \beta & \delta & \eta(s_i) \\ \circ & \times & \circ & -2 & 0 & 0 & \gamma - 2 \\ \circ & \times & \times & 0 & 0 & 0 & \gamma \\ \circ & \circ & \circ & 0 & 0 & 0 & \gamma \\ \times & \times & \circ & -2 & 0 & 2 & \gamma \\ \circ & \circ & \times & 2 & 0 & 0 & \gamma + 2 \\ \times & \circ & \circ & 0 & 0 & 2 & \gamma + 2 \\ \times & \times & \times & 0 & 0 & 2 & \gamma + 2 \\ \times & \circ & \times & 2 & 0 & 2 & \gamma + 4 \end{bmatrix}.$$

(c) If $s_i = 2d'$, then $\beta = 0$, $\gamma = -1$:

$$\begin{bmatrix} s_{i+1} - c & s_i - d & s_{i+1} - c' & \alpha & \beta & \delta & \eta(s_i) \\ \circ & \times & \circ & -2 & 0 & 0 & -3 \\ \circ & \times & \times & 0 & 0 & 0 & -1 \\ \circ & \circ & \circ & 0 & 0 & 0 & -1 \\ \times & \times & \circ & -2 & 0 & 2 & -1 \\ \circ & \circ & \times & 2 & 0 & 0 & 1 \\ \times & \circ & \circ & 0 & 0 & 2 & 1 \\ \times & \times & \times & 0 & 0 & 2 & 1 \\ \times & \circ & \times & 2 & 0 & 2 & 3 \end{bmatrix}.$$

(d) If $s_i \in [2d' + 1, c' + d - 1]$, then $\beta \in \{0, 1\}$, $\gamma = 0$:

$\nu(s_{i+1}) < \nu(s_i)$ if and only if the following row is satisfied

$$\begin{bmatrix} s_{i+1} - c & s_i - d & s_{i+1} - c' \\ \circ & \times & \circ \end{bmatrix}.$$

(e) If $s_i \in [c' + d, 2d]$, then $\beta = -1$, $\gamma = 0$, $s_i - d \in S \setminus S'$, $s_{i+1} - c' \notin S'$, then $\nu(s_{i+1}) < \nu(s_i) \iff s_i - \ell - d \notin S$.

The next theorem collects the results [6, Th. 4.1, Th.4.2, Th. 4.4] with some upgrades: statement (1) improves [6, Th.4.2], the last part of (5) is new.

Theorem 3.4 *With Setting 2.1, the following inequalities hold.*

(0) If $\tilde{s} \geq 2d' - d$, then $s_m \leq \tilde{s} + d$;
if $\tilde{s} < 2d' - d$, then $s_m \leq 2d'$. More precisely

(1) If $\tilde{s} \geq d' + c' - d$, then $s_m = \tilde{s} + d$.

(2) If $\tilde{s} = d' + c' - d - 1$, then $s_m \leq \tilde{s} + d - 1$.

(3) If $2d' - d < \tilde{s} < d' + c' - d - 1$, let

$$U := \{\sigma \in [2d' + 1 - d, \tilde{s}] \cap S \mid \sigma - \ell \notin S, \sigma + d + 1 - c' \notin S\}:$$

(a) if $U \neq \emptyset$, then $s_m = d + \max U$,

in particular $s_m = \tilde{s} + d \iff \tilde{s} + d + 1 - c' \notin S$,

(b) if $U = \emptyset$, then $s_m \leq 2d'$.

(4) If $\tilde{s} = 2d' - d$, then $s_m = \tilde{s} + d$.

(5) If $\tilde{s} < 2d' - d$, then $s_m \leq 2d'$, more precisely:

$$s_m = 2d' \iff 2d' \text{ satisfies either row 3 or row 4 of Table 3.3 (c).}$$

In the case $\tilde{s} + d + 1 - c' \notin S$:

- (a) if $2d' - d - 2 \leq \tilde{s} \leq 2d' - d - 1$, then $\tilde{s} + d \leq s_m \leq 2d'$
- (b) if $\tilde{s} = 2d' - d - j$, $j = 3, 4$ and $\{d' - j, \dots, d' - 1\} \cap S \neq \{d' - j + 1\}$ then $s_m \geq \tilde{s} + d$.

Proof. (0) is proved in [6, (4.4.1),(4.4.3)].

Now recall that $\tilde{s} \geq c - e$ (2.6.2), hence $\tilde{s} + d + 1 \geq c + 1 \in S$; further in cases (1) and (2) $\tilde{s} + d + 1 - c' \geq d'$, hence (1) and (2) follow by (0) and by Tables 3.3 (d), (e).

The cases (3) and (4) follow easily by Tables 3.3 (d) and (c).

(5) We have $s_m \leq 2d'$ by (0); further $2d'$ cannot satisfy the first two rows of Table 3.3 (c) since $\tilde{s} < 2d' - d$.

By a direct computation we can see that we always have $\gamma(2d' - j) \leq 1$, for $j \leq 2$, while for $j = 3, 4$ $\gamma(2d' - j) \leq 1 \iff \{d' - j, \dots, d' - 1\} \cap S \neq \{d' - j + 1\}$. Now (a) and (b) follow because $\nu(\tilde{s} + d) > \nu(\tilde{s} + d + 1)$ by Table 3.3.(b). \diamond

The following conjecture gives a lower bound for s_m , it is justified by calculations in very many examples. We are able to prove that it holds in many cases.

Conjecture 3.5 For every semigroup the inequality $s_m \geq c + d - e$ holds.

First we note that (3.5) holds in the following general cases:

Proposition 3.6 Assume $\left[\begin{array}{l} \text{either } s_m \geq \tilde{s} + d \\ \text{or } s_m \geq 2d' \text{ and } \tilde{s} < d' \end{array} \right]$. Then $s_m \geq c + d - e$.

In particular if either $\tilde{s} \geq c' + d' - d$ or $\tilde{s} + d = 2d'$, then $s_m \geq c + d - e$.

Proof. The first statement follows by (2.6.2): in fact $\tilde{s} \geq c - e$.

If $\tilde{s} < d'$, $s_m \geq 2d'$, we have $d' \geq c - e + \ell$ (2.6.1b) and $d - d' \leq \ell$ (2.8.1). Hence $s_m \geq 2d' \geq d' + c - e + \ell \geq c + d - e$. Now the particular cases follow by (3.4.1,4). \diamond

Corollary 3.7 (1) If $s_m > 2d'$, then $s_m - d \in S$.

(2) If $\tilde{s} = d' - 1$, then $s_m = \tilde{s} + d \iff c' \neq d$.

Proof. (1) follows by Table 3.3.(d).

(2) If $c' = d$, then $s_m \neq \tilde{s} + d$, by (3.4.2).

If $c' \neq d$ and $\tilde{s} = d' - 1$, we have $\tilde{s} \geq d' + c' - d$ then apply (3.4.1). \diamond FFFFFFFFFFFFFFFFFFFF

Proposition 3.8 Assume $\tilde{s} \leq d' - 2$ and $[\tilde{s} + 2, d'] \cap \mathbb{N} \subseteq S$. Then $s_m \leq \tilde{s} + d$:

(1) if $2d' - d < \tilde{s} < d' + c' - d$, then $s_m \left[\begin{array}{l} = \tilde{s} + d \iff \tilde{s} + 1 \notin S \text{ and } c' = d \\ \leq \tilde{s} + d - 1, \text{ otherwise.} \end{array} \right.$

(2) if $\tilde{s} \leq 2d' - d$, then $s_m = \tilde{s} + d$.

Proof. In case (1), by applying Theorem 3.4 we see that $s_m \leq \tilde{s} + d$; further $s_m = \tilde{s} + d \iff \tilde{s} + d + 1 - c' \notin S$. Since $\tilde{s} + 1 \leq \tilde{s} + d + 1 - c' \leq d'$ by the assumptions, we see that $s_m = \tilde{s} + d \iff c' = d$ and $\tilde{s} + 1 \notin S$.

In case (2), by Theorem 3.4 we have $s_m \leq 2d'$.

Now let $\tilde{s} + d + 1 \leq s \leq 2d'$. Then by the assumptions we get

$$\left\{ \begin{array}{l} \tilde{s} + 2 \leq s + 1 - c' \leq s - d' - 1 < d' \\ s + 1 \in S \text{ and } s + 1 - c' \in S' \\ \{s - d', \dots, d'\} \subseteq S \text{ (hence } \gamma(s) = -1 \text{ (3.3.2))} \\ s - \ell - d \in S \text{ (by (2.4.1)).} \end{array} \right.$$

From Tables 3.3 (b) - (c) we conclude that $s_m < s$ and also that $s_m = \tilde{s} + d$; in fact $\tilde{s} \in S$, $\tilde{s} - \ell \notin S$, further $\tilde{s} + d - d' \geq \tilde{s} + 2$, because $d - d' \geq 2$, therefore $\gamma(\tilde{s} + d) = -1$ by the assumptions and by (3.3.2). \diamond

Remark 3.9 (1) Both situations of (3.8.1) above can happen, even for $\ell = 3$ (see the following (4.7)):

(A) If $\ell = 3$, $t = 5$, $d' = d - 3$, $c' = d - 1$, (4.7.case A) we have $s_m < \tilde{s} + d$.

(B) If $\ell = 3$, $t = 5$, $d' = d - 3$, $d - 4 \notin S$, $c' = d$, (4.7.case B) we have $s_m = \tilde{s} + d$.

(2) Assume $2d' - d < \tilde{s} \leq d' + c' - d - 1$ and $[d' - \ell + 2, d'] \cap \mathbb{N} \subseteq S$; then the set U of (3.4.3) is empty. In fact for each $s \in S$, such that $2d' + 1 \leq s \leq \tilde{s} + d$, we have $s + 1 \in S$, and by (2.8.1), $d' - \ell + 2 \leq 2d' + 2 - d \leq s + 1 - c' \leq d'$, therefore $s + 1 - c' \in S'$.

(3) If $s_m < 2d' \leq \tilde{s} + d$, then $\begin{cases} (a) \tilde{s} + d + 1 - c' \in S \\ (b) \tilde{s} + d + 1 - c' - \ell \in S \\ (c) \{2d' - d - \ell, 2d' + 1 - c'\} \cap S \neq \emptyset. \end{cases}$

(3.a) holds by (3.4.3); in fact the assumptions imply $U = \emptyset$ because $\tilde{s} - \ell \notin S$.

(3.b) is clear by (3.a) and by (2.4.1), since $\tilde{s} < \tilde{s} + d + 1 - c' < d$.

(3.c) follows by Table 3.3 (c).

(4) The assumption $s_m > 2d'$ in (3.7.1) is necessary: for instance if $S = \{0, 20_e, 21, 26, 27_{d'}, 32_d, 39_c \rightarrow\}$ we have $s_m = 2d'$, with $2d' - d \notin S$ (we deduce $s_m = 2d'$ by Table 3.3 (c)).

Proposition 3.10 *If $\tilde{s} < d'$ and $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$, let $h = d - c'$, $q = d - d'$. Then*

(1) $[\tilde{s} - \ell + 1, d'] \cap \mathbb{N} \subseteq S$ and $e \geq 2\ell + t$.

(2) If $2d' - d < \tilde{s} \leq d' + c' - d - 1$, we have

(a) $q + h + 1 \leq t < 2q$ ($\leq 2\ell$),

(b) For $s \in [\tilde{s} + d' - \ell + 1, 2d'] \cap S$, we have $\gamma(s) = -1$.

(c) We have $s_m \leq 2d'$.

(d) Let $W := [\tilde{s} - 2\ell + 1, 2d' - \ell - d] \cap \mathbb{N} \setminus S$.

If $W \neq \emptyset$, let $h_0 := \max W$, then $s_m \geq h_0 + \ell + d$.

(e) $s_m < \tilde{s} + d - \ell + 1 \iff [\tilde{s} - 2\ell + 1, 2d' - d - \ell] \cap \mathbb{N} \subseteq S$,

$s_m < \tilde{s} + d - \ell + 1 \implies e \geq 3\ell + t$.

(f) If $[\tilde{s} - 2\ell + 1, 2d' - d - \ell] \cap \mathbb{N} \subseteq S$, then $s_m \geq \tilde{s} + d' - \ell + 1$.

Proof. (1) By the assumptions and by (2.4.1) we have $[\tilde{s} - \ell + 1, d'] \cap \mathbb{N} \subseteq S$; the inequality $e \geq 2\ell + t$ follows by (2.6.3)

(2) Statement (a) is immediate by the assumption $2d' - d < \tilde{s} \leq d' + c' - d - 1$.

(b) follows by (1) and by (3.3.2).

(c) By (3.4.3) we know that $s_m \leq \tilde{s} + d$. For each $2d' < s \leq \tilde{s} + d$ we have $d' - \ell < 2d' + 1 - c' < s + 1 - c' \leq \tilde{s} + d + 1 - c' \leq d' + c' - 1 + 1 - c' = d'$. Therefore $s + 1 - c' \in S$ and $s_m \leq 2d'$ by (3.4.3b).

(d) and (e) Note that $s \in [\tilde{s} + d' - \ell + 1, 2d'] \cap S \implies \tilde{s} - \ell + 1 \leq s - d \leq s + 1 - c' \leq s - d' \leq d'$, hence $\{s - d, s + 1 - c'\} \subseteq S'$ and $\gamma(s) = -1$, by (b) and the assumptions. By Table 3.3 (b) we get

$$\nu(s) > \nu(s + 1) \iff s + 1 - c \notin S.$$

Then (d) follows and the equivalence (e) becomes immediate by (c) and (d), recalling that $s + 1 - c = s - \ell - d$. We get $e \geq 3\ell + t$ by (2.6.1–2), since $d - (2\ell + t - 1) \in S$ and $2\ell + t - 1 < e$ by (1).

(f) For $s \in [\tilde{s} + d' - \ell + 1, 2d' - \ell] \cap \mathbb{N}$, we have $\gamma(s) = -1$ (see (b)). If there exists $\bar{s} \in [\tilde{s} + d' - \ell + 1, 2d' - \ell] \cap \mathbb{N}$, $\bar{s} + 1 - c' \notin S$, we have $\bar{s} - d \in S'$ by the assumptions and so $s_m \geq \bar{s}$ by Table 3.3 (b); the claim follows.

Assume on the contrary that $[\tilde{s} + d' - \ell + 2 - c', 2d' - \ell + 1 - c'] \cap \mathbb{N} \subseteq S$: then

$$[\tilde{s} - 2\ell + 1, 2d' - \ell + 1 - c'] \cap \mathbb{N} \subseteq S.$$

In fact $q + h + 1 \leq t$ (3.10.1) $\implies \tilde{s} + d' - \ell + 2 - c' = d - q - \ell - t + 2 + h \leq d - 2q - \ell + 1 = 2d' - \ell + 1 - d$. We can iterate the algorithm looking for one element $\bar{s} \in [2d' - \ell + 1, 2d' - \ell + d + 1 - c'] \cap \mathbb{N}$ such that $\bar{s} + 1 - c' \notin S$. If needed we repeat the argument till we find s' such that $s' + 1 - c' \notin S$: s' surely

exists since $\tilde{s} - \ell \notin S$. \diamond

The previous results can be summarized in the following theorem.

Theorem 3.11 *Assume $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$. Then $s_m \geq c + d - e$. In particular:*

- (1) *if $2d' - d < \tilde{s} < d' + c' - d$, we have $c + d - e \leq \tilde{s} + d' - \ell + 1 \leq s_m \leq 2d'$,*
- (2) *$s_m = \tilde{s} + d$ in the remaining cases.*

Proof. (1) Since $\tilde{s} < d'$, we have $[\tilde{s} - \ell + 1, d'] \cap \mathbb{N} \subseteq S$, $e \geq 2\ell + t$ by (3.10.1). It follows that $\tilde{s} - \ell + 1 \geq c - e$ because $[c - \ell - e, c - e - 1] \cap S = \emptyset$ and $\tilde{s} \geq c - e$ by (2.6.2).

The inequalities follow by items (c),(d), (e), (f) of (3.10):

if the set W of (3.10.2d) is not empty then we see that $s_m \geq \tilde{s} + d - \ell + 1 \geq \tilde{s} + d' - \ell + 3$ by (3.10.2d), recalling that $d' \leq d - 2$.

If $W = \emptyset$, by (3.10.2f) we get $s_m \geq \tilde{s} + d' - \ell + 1$.

(2) follows by (3.4.1) and by (3.8.2). In this case $s_m \geq c + d - e$ by (3.6).

To prove $s_m \geq c + d - e$ in case (1), first assume that $W = \emptyset$: since $d' \geq d - \ell$ (2.6.3b) and $e \geq 3\ell + t$ (3.10.2f), we get $s_m \geq \tilde{s} - \ell + d' + 1 \geq \tilde{s} - \ell + d - \ell + 1 = c + d - 3\ell - t \geq c + d - e$.

If $W \neq \emptyset$, since $e \geq 2\ell + t$ we get

$$s_m \geq \tilde{s} + d - \ell + 1 = c + d - 2\ell - t \geq c + d - e,$$

further $d > d' + 1 \implies \tilde{s} + d - \ell + 1 > \tilde{s} + d' - \ell + 1$. \diamond

4 Some particular case.

In this section we shall estimate or give exactly the value of s_m in some particular case. Since $s_m = \tilde{s} + d$ for each semigroup S satisfying $\tilde{s} \geq c' + d' - d$, in this section we shall often assume $\tilde{s} < c' + d' - d$.

4.1 Relations between the order bound and the holes set H .

Let $H := [c - e, c'] \cap \mathbb{N} \setminus S$ be as in (2.1): when H is an interval we deduce the value of s_m ; if $\#H \leq 2$ and in some other situation we give a lower bound for s_m .

Proposition 4.1 (1) *If $H = \emptyset$, then $c' = c - e$ and S is acute with $s_m = \tilde{s} + d$.*

(2) *Assume that $\tilde{s} < d'$. Then the conditions*

- (a) $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$ and $e = 2\ell + t$
- (b) $H = [d' + 1, c' - 1] \cap \mathbb{N}$

are equivalent and imply: $s_m = \begin{cases} 2d' & \text{if } 2d' - d + 1 \leq \tilde{s} \leq d' + c' - d - 1 \\ \tilde{s} + d & \text{in the remaining cases.} \end{cases}$

Proof. (1) $H = \emptyset \iff c' = c - e$; then apply (2.3. 5 and 4).

(2), (a) \iff (b). (a) implies that $[\tilde{s} - \ell + 1, d'] \cap \mathbb{N} \subseteq S$, and $\tilde{s} - \ell = c - e - 1$ (3.10.1) and (2.6.4). Hence (b) holds. On the contrary, (b) implies $[c - e, d'] \cap \mathbb{N} \subseteq S$. Since $c - e \leq d' - \ell$ by (2.6.3), we get $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$, further $e = 2\ell + t$ by (2.8.3 and 4).

Now assume that (a) – (b) hold. Since $c - e \leq d' - \ell$ by (2.6.1), when $2d' - d < \tilde{s} < d' + c' - d$, we have by (2.8.1) $c - e - \ell \leq d' - \ell - (d - d') = 2d' - \ell - d < \tilde{s} - \ell \leq c - e - 1$. We obtain $s_m \geq 2d'$ by (3.10.2d) because the set W as in (3.10.2d) has $2d' - \ell - d = \max W$. Now $s_m = 2d'$ follows by (3.10.2c). For the statement in the remaining cases see (3.11.2). \diamond

Example 4.2 When $\tilde{s} > d'$ the implication (b) \implies (a) in (4.1.2) is not true in general: in fact $S := \{0, 8_e, 12_{c-e=d'}, 14_{c'}, 15, 16_d, \dots, 20_c \rightarrow\}$ has $H = \{c' - 1\} = \{13\}$, $t = 0$, $\ell = 3$, $e \neq 2\ell + t$.

Proposition 4.3 Assume $\tilde{s} < c' + d' - d$. Let $k := \min\{n \in \mathbb{N} \mid d' - n \notin S\}$, $h := d - c'$, $s := d' + c' - k - 1$. We have

- (1) $s \leq 2d' \iff c' - d' \leq k + 1 \iff d - d' \leq k + h + 1$.
- (2) If $\tilde{s} < d' - k$ and $\ell \leq k + h + 1$, then $s_m \geq s \geq c + d - e$.
- (3) If $1 \leq k < \ell$, $c' - d' \leq k + 1$ and $\{d' - \ell, \dots, d'\} \setminus \{d' - k\} \subseteq S$, then $c + d - e \leq s \leq s_m \leq 2d'$.

Proof. (1) is obvious by the assumptions.

(2) We have $[d' - k - \ell + 1, d' - \ell] \cap \mathbb{N} \subseteq S$ (2.4.1).

Now we claim that $\gamma(s) = -1$. In fact the assumption $\ell \leq k + h + 1$ implies $c' - d' = d - d' - h \leq \ell - h \leq k + 1$, and so $s \leq 2d'$; since $[s - d', d'] \cap \mathbb{N} \subseteq S$, then $\gamma(s) = -1$ (3.1.2). Since $\tilde{s} < d'$, then $d - c' \leq \ell - 2$, further $\ell \leq k + h + 1$; therefore

$$d' - k - \ell + 1 \leq s - d = d' + c' - k - 1 - d \leq d' - \ell.$$

Hence $s - d \in S$. Moreover $s + 1 - c' = d' - k \notin S$. Then $s_m \geq s$ by Table 3.3 (b).

To prove that $s \geq c + d - e$, recall that $c' - d' \leq k + 1$. Then by assumption we have $\tilde{s} < d' - k < c' - k - 1 \leq d'$. Then $c' - k - 1 \in S$ and by (2.6.3) we get $c' - k - 1 \geq c - e + \ell$. Hence

$$s = d' + c' - k - 1 \geq d - \ell + c' - k - 1 \geq c + d - e \quad (2.8.1)$$

(3) By assumption $\tilde{s} < d' - h \leq d'$, and so $[d' - h - \ell, d' - \ell] \cap \mathbb{N} \setminus \{d' - k - \ell\} \subseteq S$. Hence $[d' - h - \ell, d' - k - 1] \cap \mathbb{N} \setminus \{d' - k - \ell\} \subseteq S$. Now recalling that $k < \ell$ we get:

$$d' - h - \ell \leq s - d < d' - k - h \leq d' - k.$$

Hence $s - d \in S$: in fact $s - d \neq d' - k - \ell$ because $s - d \geq d' - k - \ell + 1$. Further $s + 1 - c' \notin S$ and $\gamma(s) = -1$ by (3.1.2) since $s \leq 2d'$ and $\{s - d', \dots, d'\} \subseteq S$. Then $s \leq s_m$ by Table 3.3 (b).

The inequality $s \geq c + d - e$ can be proved as in (2).

In order to prove the last inequality, by the assumptions on \tilde{s} and by (3.4) it suffices to consider elements $u \in [2d' + 1, c' + d' - 1] \cap \mathbb{N}$. For such an element u we have $d' \geq u + 1 - c' > s + 1 - c' = d' - k$; hence $u + 1 - c' \in S'$, then $\nu(u + 1) \geq \nu(u)$ by Table 3.3 (d). \diamond

Corollary 4.4 Suppose $\tilde{s} < c' + d' - d$ and $\#H \leq 2$. Then $s_m \geq c + d - e$.

Proof. If $\#H = 0$, we have $s_m = \tilde{s} + d$ and we are done by (4.1.1) and (3.6).

If ($\#H = 2$ and $H = \{d' + 1, c' - 1\}$), or $\#H = 1$, then either $s_m = \tilde{s} + d$, or $s_m = 2d'$ (4.1.2); now see (3.6).

Finally assume that $H = \{d' - k\} \cup \{d' + 1\}$, with $k \geq 1$. In this case we have $c' - d' = 2 \leq k + 1$, since $k \geq 1$. Hence the claim $s_m \geq c + d - e$ follows by (3.11), if $k > \ell$, and by (4.3.3) if $k < \ell$. \diamond

4.2 Case $\ell = 2$.

If $\ell = 2$, the conjecture (3.5) is true, more precisely by [6, Thm 5.5] we have:

Proposition 4.5 Assume $\ell = 2$, then $s_m \geq c + d - e$ and

- (1) $s_m = \tilde{s} + d$ if $\begin{cases} t \leq 2, \\ t = 4 \\ t \geq 5 \text{ and } d - 3 \in S. \end{cases}$
- (2) $s_m = 2d - 4$ if $\begin{cases} \text{either } t = 3 \text{ and } d - 6 \notin S \\ \text{or } t \geq 5 \text{ and } d - 3 \notin S. \end{cases}$
- (3) $s_m = 2d - 6$ if $t = 3$ and $d - 6 \in S$ (all the remaining cases).

Proof. The value of s_m is known by [6, Thm. 5.5]. Another proof can be easily deduced by (2.4.1), (3.4.2), (4.3.3), (3.8.2), Table 3.3 (d), (3.10. d, f). The inequalities $s_m \geq c + d - e$ now follow respectively by (3.6) and by (3.11.3). \diamond

4.3 Case $\ell = 3$.

If $\ell = 3$, we compute explicitly the possible values of s_m and we show that the conjecture (3.5) holds.

Notation 4.6 (1) If $s_i = 2d - k \in S$, $k \in \mathbb{N}$, let

$$M(s_i) := \{(s_h, s_j) \in S^2 \mid s_i = s_h + s_j, s_h \leq d, s_j \leq d\}.$$

Note that $M(s_i) = \{(d - x, d - y) \in S^2, \mid 0 \leq x, y \leq k, x + y = k\}$ and that for $s_i \geq c$, we have $s_{i+1} - c = d - \ell - k$; for short it will be convenient to use the following notation.

$$(*) \quad \begin{cases} \Sigma := \{z \in \mathbb{N}, \mid z \leq d, d - z \in S\} \\ (c, h) \in S \times \Sigma, h = d + c - s_i \text{ instead of the pair } (c, s_i - c) \in N(s_i). \end{cases}$$

If $\ell = 3$, then $e \geq t + \ell + 1 = t + 4$ (2.6.2). To calculate the value of s_m , we shall assume $\tilde{s} < c' + d' - d$, otherwise $s_m = \tilde{s} + d$ by (3.4.1). Then we have $t \geq 3$, $d - 3 \in S$ and $d - 3 \leq d'$, by (2.5.2) and by (2.8.31). Three cases are possible:

$$\begin{aligned} \text{Case A: } S &= \{0, e, \dots, d - 3, *, d - 1, d, ***, c = d + 4 \rightarrow\} & d' &= d - 3 \\ \text{Case B: } S &= \{0, e, \dots, d - 3, * *, d, ***, c = d + 4 \rightarrow\} & d' &= d - 3 \\ \text{Case C: } S &= \{0, e, \dots, d - 3, d - 2, *, d, ***, c = d + 4 \rightarrow\} & d' &= d - 2. \end{aligned}$$

To describe $M(2d - k)$ we shall use the notations $(*)$ fixed in (4.6) and for when necessary for an element $2d - k$ we shall list all the pairs $(x, y) \in M'(2d - k)$ and the pair $(c, \ell + k + 1) \in S \times \Sigma$ (the pairs underlined $(\underline{\quad}, \underline{\quad})$ not necessarily belong to Σ^2).

Proposition 4.7 Assume $\ell = 3$. Then $s_m \geq c + d - e$. More precisely the values of s_m can be computed as follows.

$$\text{Case A. We have: } s_m = \begin{cases} \tilde{s} + d & \text{if either } t \in [0, 7] \setminus \{5\} \text{ or } (t \geq 8, d - 5 \in S) \\ 2d - 7 & \text{if } t \geq 8, d - 5 \notin S \\ & \text{if } t = 5: \\ 2d - 6 & \text{if } d - 9 \notin S \\ 2d - 7 & \text{if } d - 9 \in S, d - 10 \notin S \\ 2d - 9 & \text{if } \{d - 9, d - 10\} \subseteq S, d - 12 \notin S \\ 2d - 10 & \text{if } \{d - 9, d - 10, d - 12\} \subseteq S \end{cases}$$

Proof. $S = \{0, e, \dots, d - 3, *, d - 1, d, ***, c = d + 4 \rightarrow\}$, with $e \geq \ell + t + 1 = t + 4$ (2.6.1).

First we have $s_m = \tilde{s} + d$ if $t \leq 2d - c' - d' = 4$ and $s_m < \tilde{s} + d$, if $t = 5$ by (3.4.1 and 2). Hence we can assume $t \geq 5$, so that $d - 4 = d - 1 - \ell \in S$, $d - 3 - \ell = d - 6 \in S$, i.e., $\{0, 1, 3, 4, 6\} \subseteq \Sigma$.

If $t = 5$ we have $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$ and $2d' = 2d - 6 < \tilde{s} + d = d' + c' - 1$. We obtain that $2d - 10 \leq s_m \leq 2d'$, $e \geq 2\ell + t$ and $s_m \geq c + d - e$ by (3.11). More precisely we can verify that:

$$s_m = \begin{cases} 2d - 6 & \text{if } 9 \notin \Sigma & s_m = 2d' \geq c + d - e \\ 2d - 7 & \text{if } 9 \in \Sigma \text{ and } 10 \notin \Sigma & s_m = c + d - 11 \geq c + d - e \\ 2d - 9 & \text{if } 9 \in \Sigma, 10 \in \Sigma, 12 \notin \Sigma & s_m = c + d - 13 \geq c + d - e \\ & \text{in fact } 10 \in \Sigma \implies e \geq 14 \\ 2d - 10 & \text{if } \{9, 10, 12\} \subseteq \Sigma & s_m \geq c + d - e. \end{cases}$$

Note that in this case we have $d + d' - \ell - t + 1 = 2d - 2\ell - t + 1 = 2d - 10$ and this bound is achieved if $\{9, 10, 12\} \subseteq \Sigma$ (with $e \geq 16$). See, e.g. $S = \{0, 16, 34, 36, 37, 39, 40, \underline{41}, 42, 43, 45, 46, 50_c \rightarrow\}$.

If $t \geq 6$ we have $\tilde{s} \leq 2d' - d$ and we consider the following subcases.

If $t = 6$, then $\tilde{s} = 2d' - d = s_m$ by (3.4.4).

If $t \geq 7$ and $5 \in \Sigma$, one has $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$ and $\tilde{s} < 2d' - d$, hence $s_m = \tilde{s} + d$, by (3.11).

If $t \geq 7$ and $5 \notin \Sigma$, we know that $s_m \leq 2d'$ by (3.4.5); one can compute directly that

$$\nu(2d') < \nu(2d' + 1) \text{ (see Table 3.3 (c)) and that } \nu(2d - 7) > \nu(2d - 6),$$

hence $s_m = 2d - 7 = 2d' - 1$. Since $d - 6 = d' - \ell \in S$, we get $e \geq 2\ell + 1 + 6 = 13$ (2.6.3), and so $s_m \geq c + d - e + 2$.

$$\text{Case B. We have: } \left[\begin{array}{ll} \text{Case } t \leq 3: & s_m = \tilde{s} + d \\ \text{Case } t \geq 4, d-5 \notin S: & s_m = 2d-6 \\ \text{Case } t \geq 4, d-5 \in S, d-4 \in S: & \\ \quad \text{if } t \in \{4, 5\}, & s_m \in [2d-9, 2d-6] \\ \quad \text{if } t \geq 6, & s_m = \tilde{s} + d \\ \text{Case } t \geq 5, d-5 \in S, d-4 \notin S: & \\ \quad \text{if } t \in \{5, 6, 8\}, & s_m = \tilde{s} + d \\ \quad \text{if } t = 7, & s_m \in \{2d-8, 2d-11\} \\ \quad \text{if } t \geq 9, d-7 \notin S, & s_m = 2d-8 \\ \quad \text{if } t = 9, d-7 \in S, & s_m \in \{2d-10, 2d-11, 2d-13\} \\ \quad \text{if } t \geq 10, d-7 \in S, & s_m = \tilde{s} + d. \end{array} \right.$$

Proof. $S = \{0, e, \dots, d-3, * * , d, ***, c = d+4 \rightarrow\}$, $e \geq t+4$.

As in case A we can see that $s_m = \tilde{s} + d$ if $t \leq 3$ and $s_m < \tilde{s} + d$ for $t = 4$. Suppose $t \geq 4$. Then $\{0, 3, 6\} \subseteq \Sigma$. We deduce the statement by means of the following table:

$$\left[\begin{array}{ll} 2d-3 & (0, 3) \\ 2d-4 & (0, 4) \quad (c, 8) \\ 2d-5 & (0, 5) \quad (c, 9) \\ 2d-6 & (0, 6)(3, 3) \quad (c, 10). \end{array} \right.$$

If $4 \in \Sigma$, $5 \in \Sigma$, we have $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$ and by applying (3.11) we get $s_m \geq c + d - e$. More precisely, one can easily verify that for $t \in \{4, 5\}$ we have $2d-9 \leq s_m \leq 2d-6$, for $t \geq 6$ we have $\tilde{s} \leq 2d' - d$, then by (2.4.1) and by (3.8.2) we get $s_m = \tilde{s} + d$.

If $5 \notin \Sigma$, we have $s_m = 2d-6$.

If $4 \notin \Sigma$, $5 \in \Sigma$:

$$\text{we have } s_m = \left[\begin{array}{ll} 2d-5 & \iff t = 5 \\ 2d-6 & \iff t = 6 \text{ (} 8 \in \Sigma, 9 \notin \Sigma \text{)}; \end{array} \right.$$

the remaining cases to consider satisfy $\{0, 3, 5, 6, 8, 9\} \subseteq \Sigma, 4 \notin \Sigma$, with $t \geq 7$, $s_m < 2d-6$:

$$\left[\begin{array}{ll} 2d-7 & (0, 7) \quad (c, 11) \\ 2d-8 & (0, 8)(3, 5) \quad (c, 12) \quad s_m = 2d-8 \iff \left[\begin{array}{l} 7 \notin \Sigma \text{ or} \\ 7 \in \Sigma, 11 \notin \Sigma \end{array} \right. (\implies 7 \leq t \leq 8) \\ & \text{otherwise } 7, 11 \in \Sigma: \{0, 3, 5, 6, 7, 8, 9, 11\} \subseteq \Sigma, 4 \notin \Sigma \\ 2d-9 & (0, 9)(3, 6) \quad (c, 13) \\ 2d-10 & (0, 10)(3, 7)(5, 5) \quad (c, 14) \quad s_m = 2d-10 \iff 10 \in \Sigma, 13 \notin \Sigma (\implies 9 \leq t \leq 10) \\ & \text{otherwise} \quad \left[\begin{array}{l} \text{either } (\alpha) \ 10, 13 \in \Sigma \\ \text{or } (\beta) \ 10 \notin \Sigma (t = 7) \end{array} \right. \end{array} \right.$$

Case (α) : $\{0, 3, 5, 6, 7, 8, 9, 10, 11, 13\} \subseteq \Sigma, 4 \notin \Sigma$ ($\implies t \geq 9$):

$$\left[\begin{array}{ll} 2d-10 & (0, 10)(3, 7)(5, 5) \quad (c, 14) \\ 2d-11 & (0, 11)(3, 8)(5, 6) \quad (c, 15) \quad s_m = 2d-11 \iff 14 \notin \Sigma \\ & (\implies t = 9 \text{ if } 12 \notin \Sigma, t = 11 \text{ if } 12 \in \Sigma) \\ & \text{otherwise } 14 \in \Sigma: \{0, 3, 5 \leftrightarrow 11, 13\} \subseteq \Sigma, 4 \notin \Sigma \\ 2d-12 & (0, 12)(3, 9)(5, 7)(6, 6) \quad (c, 16) \quad s_m = 2d-12 \iff t = 12 \\ & \text{otherwise} \quad \left[\begin{array}{l} \text{either } (\alpha 1) \ 12 \notin \Sigma (\iff t = 9) \\ \text{or } (\alpha 2) \ 12 \in \Sigma, 15 \in \Sigma \end{array} \right. \\ 2d-13 & (0, 13)(3, 10)(5, 8)(6, 7) \quad (c, 17) \quad s_m = 2d-13 \quad \left[\begin{array}{l} \text{in case } (\alpha 1) \\ \text{in case } (\alpha 2) \iff t = 13 \end{array} \right. \\ & \text{otherwise } 16 \in \Sigma \quad \{0, 3, 5 \leftrightarrow 16\} \subseteq \Sigma, 4 \notin \Sigma: \\ 2d-14 & (0, 14)(3, 11)(5, 9)(6, 8)(7, 7) \quad (c, 18) \quad s_m = 2d-14 \iff t = 14 \\ & \text{otherwise } 17 \in \Sigma, \quad \{0, 3, 5 \leftrightarrow 15, 16, 17\} \subseteq \Sigma, 4 \notin \Sigma \dots \end{array} \right.$$

Clearly in cases $(\alpha 2)$, for each $t \geq 13$ we get $s_m = \tilde{s} + d$.

Case (β): $\{0, 3, 5, 6, 7, 8, 9, 11, \} \subseteq \Sigma$, $4, 10 \notin \Sigma$ ($t = 7$):

$$[2d - 11 \quad (0, 11)(3, 8)(5, 6) \quad \underline{(c, 15)} \quad s_m = 2d - 11.$$

$$\text{Case C. We have: } s_m = \begin{cases} \text{if } t = 3: \\ 2d - 4 & \text{if } d - 4 \in S, d - 7 \notin S \\ 2d - 5 & \text{if } (\{d - 4, d - 7\} \subseteq S, d - 8 \notin S) \text{ or } (d - 4 \notin S) \\ 2d - 7 & \text{if } \{d - 4, d - 7, d - 8\} \subseteq S \\ \text{if } t \geq 4: \\ \tilde{s} + d & \text{if } d - 4 \in S \\ 2d - 5 & \text{if } d - 4 \notin S. \end{cases}$$

Proof. $S = \{0, e, \dots, d - 3, d - 2, *, d, ***, c = d + 4 \rightarrow\}$, $d' = d - 2$.

As above we see that $t \geq 3$, $d - 3, d - 5 \in S$ (i.e. $\{0, 2, 3, 5\} \subseteq \Sigma$), $e \geq 7$. Consider the table:

$$\begin{array}{l} \left[\begin{array}{ll} 2d - 2 & (0, 2) \\ 2d - 3 & (0, 3) \\ 2d - 4 & \underline{(0, 4)}(2, 2) \quad \underline{(c, 7)} \\ 2d - 5 & (0, 5)(2, 3) \quad \underline{(c, 8)} \\ \text{if } \{4, 7, 8\} \subseteq \Sigma & \text{then } e \geq 12 \quad (d - 8 + e \geq c) \\ 2d - 6 & \underline{(0, 6)}(2, 4)(3, 3) \quad \underline{(c, 9)} \\ 2d - 7 & \underline{(0, 7)}(2, 5)(3, 4) \quad \underline{(c, 10)} \end{array} \right. \quad s_m = \begin{array}{l} 2d - 4 \quad \text{if } 4 \in \Sigma, 7 \notin \Sigma \\ 2d - 5 \quad \text{if } \left[\begin{array}{l} 4 \notin \Sigma \text{ or} \\ 4, 7 \in \Sigma, 8 \notin \Sigma. \end{array} \right. \\ \end{array} \\ \left. \begin{array}{l} \underline{(c, 8)} \\ \underline{(c, 9)} \\ \underline{(c, 10)} \\ \underline{(c, 11)} \end{array} \right. \end{array}$$

If $t = 3$, then $6 \notin \Sigma$: we get $s_m = 2d - 7$.

If $t \geq 4$ and $4 \in \Sigma$, we have $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$ and $\tilde{s} \leq 2d' - d$. Then $s_m = \tilde{s} + d \geq c + d - e$ by (3.11).

If $4 \notin S$, by the above table we deduce that $s_m = 2d - 5$. \diamond

4.4 Semigroups with CM type $\tau \leq 7$.

As a consequence of the above results we obtain lower bounds or the exact value of s_m for semigroups with small Cohen-Maculay type. First, in the next lemma we collect well-known or easy relations among the CM type τ of S and the other invariants.

Lemma 4.8 *Let τ be the CM-type of the semigroup S as in (2.1).*

(1) $\#H + \ell \leq \tau \leq e - 1$

(2) *Assume $\tau = \ell$, then $H = \emptyset$ and the following conditions are equivalent:*

(a) $\ell = e - 1$

(b) $\tau = \ell$, $c' = d$.

(c) $d = c - e$.

(d) $S = \{0, e, 2e, \dots, ke \rightarrow\}$.

(3) *If $c' > c - e$, then $\tau \geq \ell + 1$ and $\tau = \ell + 1 \implies H = \{c' - 1\}$.*

(4) *Assume $\tilde{s} \leq d'$ and $\tau = \ell + 1$. Then $\begin{cases} e \in \{2\ell + t - 1, 2\ell + t\}, & \text{if } \tilde{s} = d' \\ e = 2\ell + t, & \text{if } \tilde{s} < d'. \end{cases}$*

Proof. (1) Clearly every gap $h \geq c - e$ belongs to $S(1) \setminus S$, in particular $\{d + 1, \dots, d + \ell\} \cup H \subseteq S(1) \setminus S$. The inequality $\tau \leq e - 1$ is well-known.

(2) (3) are almost immediate.

(4). We have $\#H \leq 1$ by (1), $\tilde{s} - \ell < \tilde{s} \leq d' < c' - 1$. If $\#H = 0$, then $c' = c - e$ (4.1) and so $e = 2\ell + t$ (2.7.1). If $\#H = 1 \implies \tilde{s} - \ell \notin H$ and $d' = c' - 2$: it follows $e \leq 2\ell + t$ by (2.6.4). Now apply (2.6.3) and (2.8.2): if $\tilde{s} < d'$, then $\tilde{s} + 1 \in S$, hence $e \geq 2\ell + t$ and so $e = 2\ell + t$. Further $\tilde{s} = d' \implies c' \geq c - e + \ell \implies e \geq c - c' + \ell = d + 2\ell - d' - 1 = 2\ell + t - 1$.

Example 4.9 (1) We recall that in general $c' = c - e$ does not imply $\tau = \ell$. For instance, let $S = \{0, 10_{e=d'}, 16_{c-e}, 17, 18, 19, 20, 21, 22, 23, 24_d, 26_c \rightarrow\}$. Then $\tau = 5 \neq \ell$.

(2) Analogously $H = \{c' - 1\}$ does not imply $\tau = \ell + 1$:

$S = \{0, 10_e, 16_{c-e}, 17, 18, 19, 20_{d'}, 22_{c'=d}, 26_c \rightarrow\}$ has $\ell = 3$, $\tau = 5$.

(3) In (4.8.4) the conditions $e = 2\ell + t$ and $\tilde{s} < d'$ are not equivalent, further when $\tilde{s} = d'$ both the cases with $\tau = \ell + 1$, $\tilde{s} = d'$ are possible. For instance

$\{0, 9_{e=c-e}, 10, 11_{d'}, 13_{c'}, 14_d, 18_c \rightarrow\}$ has $t = \ell = 3$, $e = 2\ell + t$, $\tilde{s} = d'$, $\tau = \ell + 1$;
 $\{0, 8_e, 9_{d'}, 11_{c'}, 12_d, 16_c \rightarrow\}$ has $t = \ell = 3$, $e = 2\ell + t - 1$, $\tilde{s} = d'$, $\tau = \ell + 1$.

(4) There exist semigroups with $H = \{c' - 1\}$, $\tilde{s} < d'$, $\tau = \ell + 1$ as in (4.8.4):

$S = \{0 * \dots * 11_e * \dots * 15_{d-e} * \dots * 19, 20, 21, 22, 23_{d'} * 25_{c'} * 26_d * \dots * 30_c \rightarrow\}$,
has $\ell = 3$, $t = 5$, $e = 2\ell + t$, $\tau = 4$.

Now we deduce bounds for s_m when $\tau \leq 7$.

Proposition 4.10 *For each $\tau \leq 7$ we have $s_m \geq c + d - e$. More precisely when $\tilde{s} < c' + d' - d$ we have the following results.*

(1) $\tau \leq 3$. We have: $s_m = \begin{cases} 2d - 4 & \text{if } S \text{ non-acute, } \tau = t = 3 \ (\ell = 2) \\ \tilde{s} + d & \text{in the other cases} \end{cases}$
[6, 5.9] [5, 4.13]

(2) $\tau = 4$. We have $\ell \leq 4$ and the following subcases.

If $\ell = 4 (= \tau)$, then $H = \emptyset$ (4.8.1), therefore S is acute with $s_m = \tilde{s} + d$ (2.3.4).

If $\ell \leq 3$ we are done by the previous (4.5), (4.7), (2.3.4) and (4.1) (recall $\ell = 1 \implies S$ is acute).

More precisely we get:

$$s_m = \begin{cases} 2d - 4 & \text{if } \begin{cases} \ell = 2 \text{ and either } (t = 3, d - 6 \notin S) \text{ or } (t \geq 5, d - 3 \notin S) \\ \ell = t = 3, e = 9, c' = d, d' = c' - 2, d - 4 \in S \end{cases} \\ 2d - 6 & \text{if } \begin{cases} \ell = 2, t = 3, d - 6 \in S \\ \ell = 3, t = 5, e = 11, c' = d - 1, d' = c' - 2 \end{cases} \\ \tilde{s} + d & \text{in the other cases.} \end{cases}$$

(3) $\tau = 5$. As above we know s_m in every case:

(a) If $\ell \leq 3$ we can deduce s_m by (4.5), (4.7),

(b) If $4 \leq \ell \leq 5$, then we are done by (4.1), since $\#H \leq 1$.

(4) $\tau = 6$. We can calculate s_m as follows:

(a) If $\ell \leq 3$ as in (3.a).

(b) If $5 \leq \ell \leq 6$, then we are done by (4.1), since $\#H \leq 1$.

(c) If $\ell = 4$ and $H \subseteq \{c' - 2, c' - 1\}$, we have the value of s_m by (4.1) and (4.1.2).

(d) If $\ell = 4$ and $H = \{d' - k, c' - 1\}$, with $k \geq 1$, then $d' = c' - 2$, and the bounds for s_m are given in (3.11) if $k > \ell$, and in (4.4) if $\ell > k \geq 1$ (in fact $c' - d' = 2 \leq k + 1$).

(5) $\tau = 7$. We have $\ell \leq 7$ and the following subcases.

(a) If $\ell \leq 3$, then s_m is known as in (3.a).

(b) If $6 \leq \ell \leq 7$ then $\#H \leq 1$ and we are done by (4.1).

(c) If $\ell = 5$, then $\#H \leq 2$ and we are done by (4.1), (4.4).

(d) If $\ell = 4$, then $\#H \leq 3$ and we are done if $\#H \leq 2$ by (4.1), (4.4).

If $\ell = 4$, $\#H = 3$, consider the following subcases

(i) $H = [d' + 1, c' - 1] \cap \mathbb{N}$: then s_m is given in (4.1).

(ii) $H = \{d' - k, c' - 2, c' - 1\}$, $k \geq 2$. If $k < \ell$, then s_m is given in (4.3.3). If $k \geq \ell$, apply (3.11).

(iii) $H = \{d' - 1, c' - 2, c' - 1\}$: this case cannot exist. In fact since $S = \{e, \dots, d' - 2, *, d', *, *, c', \dots, d, c \rightarrow\}$ and by the assumption $\tilde{s} < d'$, we obtain $c' - \ell \in S$ and $c' - \ell = c' - 4 = d' - 1 \notin S$, impossible.

(iv) $H = \{d' - j, d' - k, c' - 1\}$, $j > k \geq 1$, hence $S = \{e, \dots, d' - 2, \dots, d', *, c', \dots, d(\leq c' + 2), \dots, d + 5 \rightarrow\}$, with $2 = c' - d' \leq k + 1$. As in the proof of (4.3.2) for $s := d' - k + c' - 1$ we have $\gamma(s) = -1$, $s + 1 - c' \notin S$, $s - d \in S \iff s - d \neq d' - j$. Hence (Table 3.3 (b)) $s_m \geq s$ if $s - d \neq d' - j$, i.e., $d' - k + c' - 1 - d \neq d' - j$, i.e., $d - c' \neq j - k - 1$.

Four subcases: $\left[\begin{array}{l} j = k + 1 \implies s_m \geq s \text{ if } d \neq c' \\ j = k + 2 \implies s_m \geq s \text{ if } d \neq c' + 1 \\ j = k + 3 \implies s_m \geq s \text{ if } d \neq c' + 2 \\ j \geq k + 4 \implies s_m \geq s \end{array} \right.$ (since $d - c' \leq \ell - 2 = 2$, $j - k - 1 \geq 3$).

In the remaining three situations we can see that $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$, therefore s_m is given by (3.11):

- If $j = k + 1$ and $d = c'$, since $\tilde{s} < d'$, $\ell = 4$ we get $\{d' - 4, d - 4 = d' - 2, d'\} \subseteq S$. Since there are two consecutive holes, then $k \geq 5$. It follows $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$.

- If $j = k + 2$, and $d = c' + 1$, we have $c' - 4 = d' - 2$, $d' - 1 = d - 4$ and $\tilde{s} < d' - 1$ (4.3.1). Therefore $S = \{e, \dots, d' - 5, d' - 4, d' - 3, *, d' - 2, d' - 1, d', *, d' + 2 = c', d' + 3 = d, d + 5 \rightarrow\}$. We deduce $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$.

- If $j = k + 3$, and $d = c' + 2$, analogously we deduce $[d' - \ell, d'] \cap \mathbb{N} \subseteq S$. \diamond

4.5 The value of s_m for semigroups of multiplicity $e \leq 8$.

Corollary 4.11 For each semigroup S of multiplicity $e \leq 8$ we have $s_m \geq c + d - e$.

Proof. Since $\tau \leq e - 1$ the result follows by (4.10). \diamond

4.6 Almost arithmetic sequences and Suzuki curves.

Recall that a semigroup S is generated by an almost arithmetic sequence (shortly AAS) if

$$S = \langle m_0, m_1, \dots, m_{p+1}, n \rangle$$

with $m_0 \geq 2$, $m_i = m_0 + \rho i$, $\forall i = 1, \dots, p + 1$, and $\text{GCD}(\rho, m_0, n) = 1$. (The embedding dimension of S is $\text{embdim } S = p + 2$).

Proposition 4.12 Let S be an AAS semigroup of embedding dimension μ ; then $\tau \leq 2(\mu - 2)$.

Proof. It is a consequence of [8, 3.3 - 4.6 - 4.7 - 5.6 - 5.7 - 5.8 - 5.9] after suitable calculations.

Corollary 4.13 If S is an AAS semigroup with $\text{embdim } S \leq 5$ then $s_m \geq c + d - e$.

Proof. It is an immediate consequence of (4.10) and (4.12). \diamond

As another corollary we obtain the value of s_m for the Weierstrass semigroup of a Suzuki curve, that is a plane curve C defined by the equation

$$y^b - y = x^a(x^b - x), \text{ with } a = 2^n, b = 2^{2n+1}, n > 0.$$

Some applications of these curves to AG codes can be found for example in [4].

Proposition 4.14 *If S is the Weierstrass semigroup of a Suzuki Curve, then S is symmetric, therefore $s_m = \tilde{s} + d$.*

Proof. In [4, Lemma 3.1] is proved that the Weierstrass semigroup S at a rational place of the function field of C is generated as follows:

$$S = \langle b, b + a, b + \frac{b}{a}, 1 + b + \frac{b}{a} \rangle$$

We have $b = 2a^2$, with $a = 2^n$, hence $S = \langle 2a^2, 2a^2 + a, 2a^2 + 2a, 2a^2 + 2a + 1 \rangle$ Then consider the semigroup

$$S = \langle 2a^2, 2a^2 + a, 2a^2 + 2a, 2a^2 + 2a + 1 \rangle, \quad a \in \mathbb{N}.$$

If $a = 1$, then $S = \langle 2, 3 \rangle$.

If $a > 1$, then S is generated by an almost arithmetic sequence, and $\text{embedim}(S) = 4$; in fact

$$S = \langle m_0, m_1, m_2, n \rangle, \quad \text{with } m_0 = 2a^2, \quad m_1 = m_0 + a, \quad m_2 = m_1 + a, \quad n = m_2 + 1.$$

Since S is AAS, we shall compute the Apery set \mathcal{A} by means of the algorithm described in [8]:

let $p = \text{embedim}(S) - 2 = 2$ and for each $t \in \mathbb{N}$,

let q_t, r_t be the (unique) integers such that $t = pq_t + r_t (= 2q_t + r_t)$, $q_t \in \mathbb{Z}$, $r_t \in \{1, 2\}$,

let $g_t = q_t m_2 + m_{r_t}$, i.e., $g_t = \begin{cases} (q_t + 1)m_2 & \text{if } r_t = 2 \\ q_t m_2 + m_1 & \text{if } r_t = 1 \end{cases}$ (in particular $g_0 = 0$).

Then by [8] the Apery set \mathcal{A} of S is: $\{ g_t + hn \mid 0 \leq t \leq 2a - 1, 0 \leq h \leq a - 1 \}$: therefore the elements of the Apery set are the $2a^2$ entries of the following matrix

$$\begin{bmatrix} 0 & g_1 & g_2 & g_3 & \cdot & \cdot & \cdot & g_{2a-1} \\ & \parallel & \parallel & \parallel & \cdot & \cdot & \cdot & \parallel \\ & m_1 & m_2 & m_1 + m_2 & \cdot & \cdot & \cdot & (a-1)m_2 + m_1 \\ n & g_1 + n & g_2 + n & \cdot & \cdot & \cdot & \cdot & \cdot \\ 2n & g_1 + 2n & g_2 + 2n & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot \\ \cdot & \cdot \\ (a-1)n & g_1 + (a-1)n & g_2 + (a-1)n & \cdot & \cdot & \cdot & \cdot & g_{2a-1} + (a-1)n \end{bmatrix}$$

Recall that a semigroup S of multiplicity e and Apery set \mathcal{A} is

symmetric \iff for each $s_i \in \mathcal{A}$, $0 < s_i \neq s_e := \max \mathcal{A}$, there exists $s_j \in \mathcal{A}$ such that $s_i + s_j = s_e$.

In our case this condition is satisfied: in fact $s_i = \begin{cases} \alpha m_2 + h n, & h \leq a - 1, \alpha \geq 1 \quad (1) \text{ or} \\ \alpha m_2 + m_1 + h n & 0 \leq \alpha, h \leq a - 1 \quad (2) \end{cases}$

further $s_e = (a - 1)(m_2 + n) + m_1$, and so

$$s_e - s_i = \begin{cases} (a - 1 - \alpha)m_2 + m_1 + (a - 1 - h) n \in \mathcal{A} & (1) \text{ or} \\ (a - 1 - \alpha) m_2 + (a - 1 - h) n \in \mathcal{A} & (2) \end{cases}$$

Since a semigroup S is symmetric if and only if its CM-type is one, then $s_m = \tilde{s} + d$ by (4.10.1). \diamond

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